

Fault tolerance with transactions: past, present and future

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Fault tolerance

Transaction fundamentals

- What is a transaction?
- ACID properties
- Distributed transactions
- Web Services

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- Machines and software fail
 - Fundamental universal law
 - Things get better with each generation, but still statistically significant
- Failures of centralized systems difficult to handle
- Failures of distributed systems are much more difficult



Fault tolerance techniques

Replication of resources

- Increase availability
 - Probability is that a critical number of resources remain operational
 - "Guarantee" forward progress
- Tolerate programmer errors by heterogeneous implementations

Spheres of control

"Guarantee" no partial completion of work in the presence of failures



Affect of time

- Fault tolerance has always been extremely important
- Back in the 1980's many different efforts
 - Emerald, Argus, Arjuna, Camelot/Avalon, Isis, Horus etc.
 - Mostly concentrated around distributed systems
 - Centralized system as degenerate case
- 1990's saw standardization of distributed systems
 - Ansa, DCE, COM/DCOM, CORBA, J2EE





Is there still research potential?

- What we do is changing
- How we do it is changing
- Paradigm shifts occurring frequently
 - Web Services
 - Grid Computing
 - Mobile Computing
 - Large Scale Computing
- These often require new techniques for fault tolerance
 - Some research efforts in environments like these started decades ago



What is a transaction?

- Mechanistic aid to achieving correctness
- Provides an "all-or-nothing" property to work that is conducted within its scope
 - Even in the presence of failures
- Ensures that shared resources are protected from multiple users



- Atomicity
- Consistency
- Isolation
- Durability







Within the scope of a transaction

- all changes occur together OR no changes occur
- Atomicity is the responsibility of the Transaction Manager

• For example - a money transfer

- debit removes funds
- credit add funds
- no funds are lost!

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Two-phase commit

- Required when there are more than one resource managers (RM) in a transaction
- Managed by the transaction manager (TM)
- Uses a familiar, standard technique:
 - marriage ceremony Do you? I do. I now pronounce ..
- Two phase process
 - voting phase can you do it?
 - Attempt to reach a common decision
 - action phase if all vote yes, then do it.
 - Implement the decision









Handling failures

Presumed Abort Strategy

- can be stated as « when in doubt abort »
- any failure prior the commit phase lead to abort the transaction

• A coordinator or a participant can fail in two ways

- it stops running (crashes)
- it times out waiting for a message it was expecting
- A recovered coordinator or participant uses information on stable storage to guide its recovery



2PC: optimizations

• one phase commit

- no voting if transaction tree is single branch





"read-only"

- resource doesn't change any data
- can be ignored in second phase of commit



Nested transactions



 a transaction is *nested* when it executes within another transaction

nested transactions live in a tree structure

- parents
- children

implement modularity and containment



Consistency

- Transactions scope a set of operations
- Consistency can be violated within a transaction
 - Allowing a debit for an empty account
 - Debit without a credit during a Money Transfer
 - Delete old file before creating new file in a copy
- transaction must be correct according to application rules
- Begin and commit are points of consistency
- Consistency preservation is a property of a transaction, not of the TP system (unlike the A, I, and D of ACID)



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State transformations new state under construction



Commit



- Running programs concurrently on same data can create concurrency anomalies
 - the shared checking account example





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- Transaction must operate as a black box to other transactions
- Multiple programs sharing data requires concurrency control
- When using transactions
 - programs can be executed concurrently
 - BUT programs appear to execute serially



Isolation







Durability

- When a transaction commits, its results must survive failures
 - must be durably recorded prior to commit
 - system waits for disk ack before acking to user
- If a transaction rolls back, changes must be undone
 - before images recorded
 - undo processing after failure



Heuristics

- Two-phase commit protocol is blocking in order to guarantee atomicity.
 - Participants may be blocked for an indefinite period due to failures
- To break the blocking nature, prepared participants may make autonomous decisions to commit or rollback
 - Participant *must* durably record this decision in case it is eventually contacted to complete the original transaction
 - If the decision differs then the coordinator's choice then a possibly nonatomic outcome has happened: a *heuristic outcome*, with a corresponding *heuristic decision*.



Interposition

- Allows a subordinate coordinator to be created
- Interposed coordinator registers with transaction originator
 - Form tree with parent coordinator
 - Application resources register locally





Web Services and SOA

- Transactions today imply all ACID properties
- Good for "short" durations
 - Application specific
- Long-running transactions may impose constraints
 - Hours, days, months, ...
 - Retain resources for duration of transaction



Web Services transactions

- Business-to-business interactions may be complex
 - involving many parties
 - spanning many different organisations
 - potentially lasting for hours or days
- Cannot afford to lock resources on behalf of an individual indefinitely
- May need to undo only a subset of work



- Internal isolation or resources should be a decision for the service provider
 - E.g., commit early and define compensation activities
 - However, it does impact applications
 - Some users may want to know a priori what isolation policies are used

Undo can be whatever is required

- Before and after image
- Entirely new business processes



- Sometimes it may be desirable to cancel some work without affecting the remainder
 - E.g., prefer to get airline seat now even without travel insurance
- Similar to nested transactions
 - Work performed within scope of a nested transaction is provisional
 - Failure does not affect enclosing transaction
- However, nested transactions may be too restrictive
 - Relaxing isolation



Structuring transactions

- Could structure transactional applications from short-duration transactions
 - Release locks early
 - Resulting application may still be required to appear to have "ACID" properties
 - May require application specific means to restore consistency
- A transactional workflow system could be used to script the composition of these transactions



Structuring transactions





Extended transaction models

• There are a number of such models

- Sagas
- Compensations
- Epsilon Serialisability
- Versioning Schemes
- Nested top-level transactions
- Open-nested transactions
- Glued transactions
- Coloured actions



Future directions

- One size does not fit all!
- Business domains will impose different requirements on implementers
 - Essentially construct domain-specific models
 - Real-time
- The range and requirements for such extended models are not yet known
 - Do not restrict implementations because we don't know what we want yet
- Still a very active area of research and development



Any questions?

