The evolution of a transaction processing system

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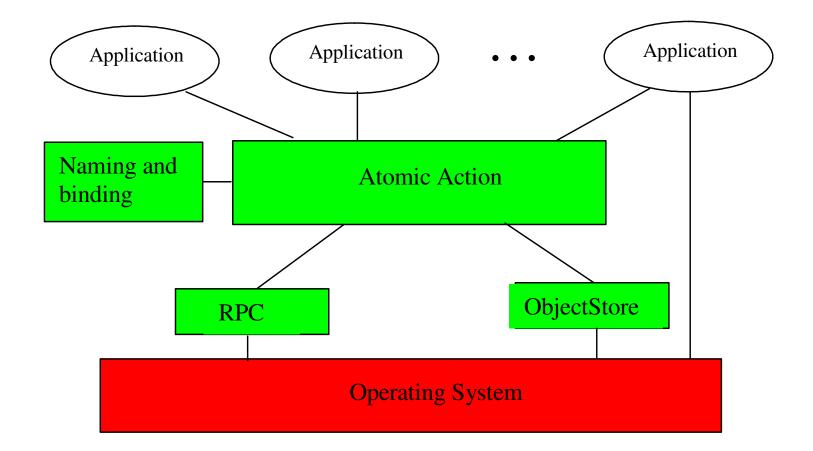
ATS background

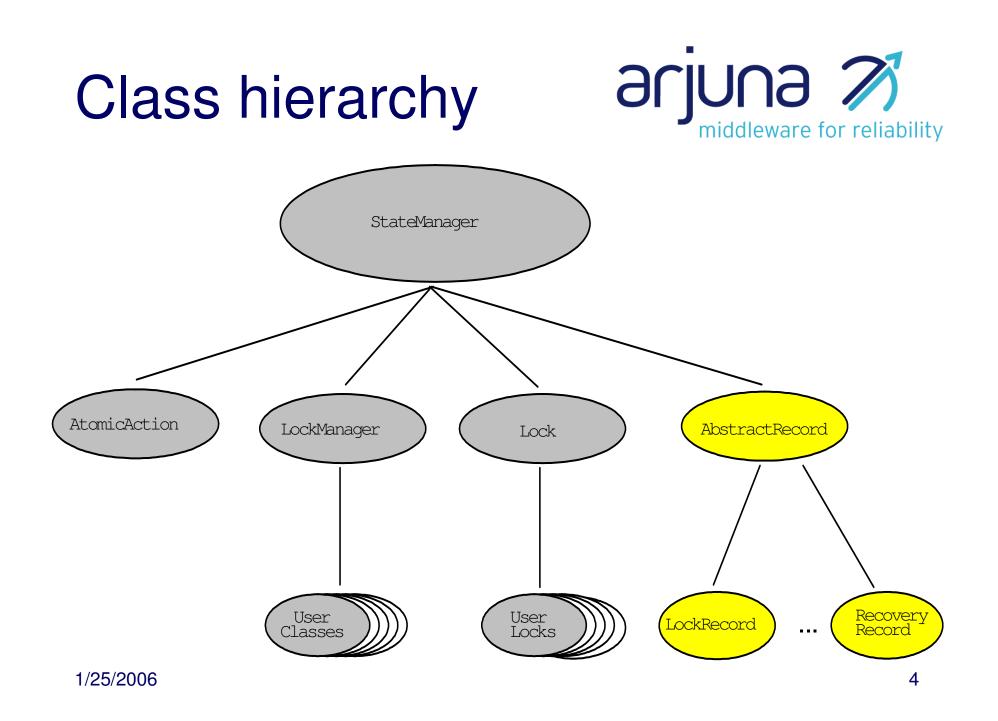


- Distributed transaction processing system
 - Began life in C++ back in 1986 at the University of Newcastle upon Tyne, England
 - Exploit object-oriented techniques
 - Pre-CORBA, DCE, COM, ...
 - Own RPC and stub-generation mechanisms
- Complete toolkit for development of faulttolerant applications
 - Persistence, concurrency control,...









AbstractRecords



- Forms the basic interface for all transaction participants
 - (nested) two-phase commit aware
 - Does not imply a specific implementation.
- Key to the longevity of Arjuna.
 - Many transaction systems then and today tie transaction participants to X/Open XA compliant resources (e.g., databases).

Basics covered



- High availability
- Standards evolution
- Performance, performance, performance!
- Support for multiple models and protocols

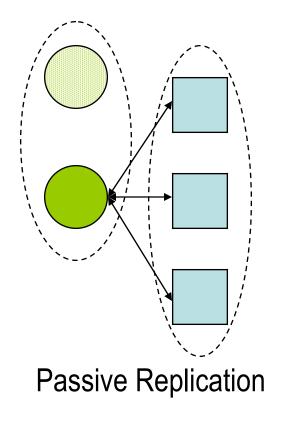
High availability

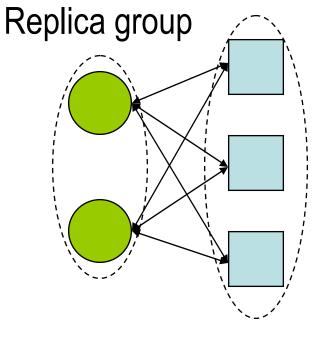


- Active replication
 - Assumes determinism
 - K+1 replicas to tolerate K failures
 - 2K+1 if network can partition
 - Group communication
 - Typically ordered delivery
- Passive replication
 - Does not require determinism
 - K+1 replicas to tolerate K failures
 - Slower fail-over time

Replication protocols







Active Replication

Student registration arjuna 7

- No money to buy
 - Nothing available to by at that time
- Must work on PC, Mac and various Unix workstations
 - 20000 students over 5 days
 - Cannot tolerate failure as student gets no money
- Campus wide
 - 10 servers, with 150 front-ends
- Network can partition

Standards evolution arjuna 7

- 1995 saw release of initial OTS specification from OMG
 - Shares many similarities with Arjuna
 - Generic two-phase participants
 - Optional support for nested transactions
 - Only a two-phase commit protocol engine
 - Persistence and concurrency control elsewhere
- Overlap in several areas
 - Naming and binding
 - RPC

Modifications



- Replace RPC and Naming/Binding modules
 - Slight modifications due to different distribution model
 - E.g., Arjuna had support for passing pointers and associated memory, CORBA IDL did not
- Transaction engine remained *unchanged*
 - Wrap OTS participants in AbstractRecords
 - Benefits from previous 10 years of testing and use

Coordinator performance



- Supported typical optimizations
 - Presumed abort, one-phase commit, read-only participants
- Also supports embedding of coordinator
 - Small footprint
 - Can run in less than 16 Meg
 - Durability and recovery are loaded on demand
 - Log structure is created on demand
 - Implementation is flexible too (no requirement for db, for example)

Coordinator performance



- Different types of participant
 - Recoverable
 - Two phase (and nesting) aware
 - Do not have any persistent state representation
 - Do not require recovery
 - Do not require (transaction) log
 - Durable recoverable
 - Have persistent state representation
 - Require recovery
 - Require (transaction) log

Multiple models and protocols



- Factor out core transaction engine
 - Essentially the same engine that began life in 1986
 - No dependency on any distribution infrastructure
 - Purely local transactions and recovery
 - Hooks for distribution are essentially the interfaces to the RPC, Naming/Binding and Crash Recovery modules
 - Participant implementations are opaque to the transaction engine
 - Context information via XML+SOAP, IIOP, ...
- Embedded within
 - HP products (HP-TS, HP-MS, HP-WST)
 - HP proof of concepts technologies (mobile devices)
 - Arjuna products (A-TS, A-MS, A-XTS)

Conclusions and lessons learnt



- Modularity helped us a lot
- AbstractRecord made it easier to customise
- Customer feedback has been extremely useful
 - "Transaction semantics are great, but relax the properties."
- Standards are good
 - But their lifetimes and impact are sometimes over hyped
- Transactions everywhere is a good idea
 - Just make them cheap to use!